252-0538-00L, Spring 2017

## Shape Modeling and Geometry Processing

#### Geometry Acquisition Meshes



February 21, 2018











Roi Poranne

**ETH** zürich













#### **Touch Probes**







### **Touch Probes**

- Physical contact with the object
- Manual or computer-guided
- Advantages:
  - Can be very precise
  - Can scan any solid surface
- Disadvantages:
  - Slow, small scale
  - Limited for fragile objects





ETH zürich



## **Optical Scanning**

- Infer the geometry from light reflectance
- Advantages:
  - Less invasive than touch
  - Fast, large scale possible
- Disadvantages:
  - Difficulty with transparent and shiny objects





## Optical scanning - active lighting Time of flight laser

- A type of laser rangefinder (LIDAR)
- Measures the time it takes the laser beam to hit the object and come back







### Optical scanning - active lighting

- Accommodates large range up to several miles (suitable for buildings, rocks)
- Lower accuracy
  - objects move while scanning





- Laser beam and camera
- Laser dot is photographed
- The location of the dot in the image allows triangulation: we get the distance to the object







- Laser beam and camera
- Laser dot is photographed
- The location of the dot in the image allows triangulation: we get the distance to the object







- Laser beam and camera
- Laser dot is photographed
- The location of the dot in the image allows triangulation: we get the distance to the object











- Very precise (tens of microns)
- Small distances (meters)







## Optical scanning - active lighting

#### Structured light



## Optical scanning - active lighting Structured light



- Very fast 2D pattern at once
  - Used in Kinect, RealSense, Face ID
- Complex distance calculation, prone to noise



# Optical scanning - passive

Stereo (x, y, z) **Right camera** focal point Epipolar line Left camera focal point  $(x_R, y_R)$ Right camera  $(x_L, y_L)$ projection plane eft camera projection plane



## Optical scanning - passive Stereo



- No need for special lighting/radiation
- Requires two (or more) cameras
  - Feature matching and triangulation



## Imaging

- Ultrasound, CT, MRI
- Discrete volume of density data
- First need to segment the desired object (contouring)









## Meshes



February 21, 2018

Boundary representations of objects





#### Meshes as Approximations of Smooth Surfaces

- Piecewise linear approximation
  - Error decreases with element size





#### Meshes as Approximations of Smooth Surfaces

- Piecewise linear approximation
  - Error decreases with element size





Meshes as Approximations of Smooth Surfaces

- Piecewise linear approximation
  - Error decreases with element size





Polygonal meshes are a good representation

 approximation arbitrary topology piecewise smooth surfaces adaptive refinement efficient rendering

Roi Porann

February 21, 2018

Hzürich

## Polyline/Polygon





 A finite set *M* of closed, simple polygons *f<sub>j</sub>* is a polygonal mesh







- A finite set *M* of closed, simple polygons *f<sub>j</sub>* is a polygonal mesh
- Every edge belongs to at least one polygon





- A finite set *M* of closed, simple polygons *f<sub>j</sub>* is a polygonal mesh
- Every edge belongs to at least one polygon
- Each f<sub>j</sub> defines a face of the polygonal mesh





- A finite set *M* of closed, simple polygons *f<sub>j</sub>* is a polygonal mesh
- Every edge belongs to at least one polygon
- Each f<sub>j</sub> defines a face of the polygonal mesh





- A finite set *M* of closed, simple polygons *f<sub>j</sub>* is a polygonal mesh
- Every edge belongs to at least one polygon
- Each f<sub>j</sub> defines a face of the polygonal mesh





- A finite set *M* of closed, simple polygons *f<sub>j</sub>* is a polygonal mesh
- Every edge belongs to at least one polygon
- Each f<sub>j</sub> defines a face of the polygonal mesh



#### Vertex **degree** or **valence** = number of incident edges




#### Polygonal Mesh

#### Vertex **degree** or **valence** = number of incident edges





# Polygonal Mesh



**Boundary:** the set of all edges that belong to only one face

- Either empty or forms closed loops
- If empty, then the polygonal mesh is closed



# **Triangle Meshes**

- Connectivity: vertices, edges, triangles
  - AKA Topology, Combinatorics
- Geometry: vertex positions
  - $V = \{v_1, \dots, v_n\}$   $E = \{e_1, \dots, e_k\}, \quad e_i \in V \times V$   $F = \{f_1, \dots, f_m\}, \quad f_i \in V \times V \times V$

$$P = \{\mathbf{p}_1, \dots, \mathbf{p}_n\}, \quad \mathbf{p}_i \in \mathbb{R}^3$$



#### Manifolds

A surface is a closed **2D manifold** if it is everywhere locally **homeomorphic** to a disk **"Looks like"** 



#### Manifolds

#### Boundary points are homeomorphic to a half-disks



# Quiz

 For each case, decide if it is a 2-manifold (possibly with boundary) or not. If not, explain why not.



#### Case 1Case 2Case 3Case 4Case 5



# Quiz

 For each case, decide if it is a 2-manifold (possibly with boundary) or not. If not, explain why not.





## Manifold meshes

- In a manifold mesh, at most 2 faces share an edge
  - Boundary edges: have one incident face
  - Inner edges have two incident faces
- A manifold vertex has 1 connected **ring** of faces around it, or 1 connected half-ring (boundary)



#### Manifold meshes

- If closed and not intersecting, a manifold divides the space into inside and outside
- A closed manifold polygonal mesh is called polyhedron





#### Orientation

- Every face of a polygonal mesh is orientable
  - Clockwise vs. counterclockwise order of face vertices
  - Defines sign/direction of the surface normal





#### Orientation

• Consistent orientation of neighboring faces:





# Orientability

- A polygonal mesh is orientable, if the incident faces to every edge can be consistently oriented
  - If the faces are consistently oriented for every edge, the mesh is oriented

#### Notes

- Every non-orientable closed mesh embedded in  $\mathbb{R}^3$  intersects itself
- The surface of a polyhedron is always orientable





# **Global Topology of Meshes**

- Genus:  $\frac{1}{2}$  × the maximal number of closed paths that do not disconnect the graph.
  - Informally, the number of handles ("donut holes").





# Global Topology of Meshes

- **Genus:**  $\frac{1}{2}$  × the maximal number of closed paths that do not disconnect the graph.
  - Informally, the number of handles ("donut holes").





# Global Topology of Meshes

• Genus: ½ × the maximal number of closed paths that do not disconnect the graph.





• Theorem (Euler): The sum

$$\chi(M) = v - e + f$$

is **constant** for a given surface topology, no matter which (manifold) mesh we choose.

- v = number of vertices
- e = number of edges
- f = number of faces



• For orientable meshes:

$$v - e + f = 2(c - g) - b = \chi(M)$$

- c = number of connected components
- g = genus
- b = number of boundary loops

$$\chi(\bigcirc) = 2 \quad \chi(\bigcirc) = 0$$



Count #edges and #faces in a closed triangle mesh:

- Ratio of edges to faces:  $2e = 3f \Rightarrow e = \frac{3}{2}f$ 
  - each edge belongs to exactly 2 triangles
  - each triangle has exactly 3 edges



54



Count #edges and #faces in a closed triangle mesh:

- Ratio of edges to faces:  $2e = 3f \Rightarrow e = \frac{3}{2}f$ 
  - each edge belongs to exactly 2 triangles
  - each triangle has exactly 3 edges
- Ratio of vertices to faces:  $f \sim 2v$

• 
$$2 = v - e + f = v - \frac{3}{2}f + f$$
  
•  $2 + \frac{f}{2} = v$ 



Count #edges and #faces in a closed triangle mesh:

- Ratio of edges to faces:  $2e = 3f \Rightarrow e = \frac{3}{2}f$ 
  - each edge belongs to exactly 2 triangles
  - each triangle has exactly 3 edges
- Ratio of vertices to faces:  $f \sim 2v$

• 
$$2 = v - e + f = v - \frac{3}{2}f + f$$

$$2 + \frac{f}{2} = v$$

- Ratio of edges to vertices:  $e \sim 3v$
- Average degree of a vertex: 6



# Regularity

- Triangle mesh: average valence = 6
- Quad mesh: average valence = 4



- **Regular mesh:** all faces have the same number of edges and all vertex degrees are equal
- Quasi-regular mesh:
  - Almost all of vertices have degree 6 (4).



#### Quad vs. Triangle Meshes

• Edge loops are great for editing





# Triangulation



- Every face is a triangle
- Simplifies data structures
- Simplifies rendering
- Simplifies algorithms
- Each face planar and convex
- Any polygon can be triangulated



# Triangulation



- Every face is a triangle
- Simplifies data structures
- Simplifies rendering
- Simplifies algorithms
- Each face planar and convex
- Any polygon can be triangulated



#### Data Structures



- What should be stored?
  - Geometry: 3D coordinates
  - Connectivity
    - Adjacency relationships
  - Attributes
    - Normal, color, texture coordinates
    - Per vertex, face, edge



#### Data Structures



- What should be supported?
  - Rendering
  - Geometry queries
    - What are the vertices of face #2?
    - Is vertex A adjacent to vertex H?
    - Which faces are adjacent to face #1?
  - Modifications
    - Remove/add a vertex/face
    - Vertex split, edge collapse



### Data Structures



- How good is a data structure?
  - Time to construct
  - Time to answer a query
  - Time to perform an operation
  - Space complexity
  - Redundancy
- Criteria for design
  - Expected number of vertices
  - Available memory
  - Required operations
  - Distribution of operations



# Triangle List

- STL format (used in CAD)
- Storage
  - Face: 3 positions
- No connectivity information

Triangles				
0	x0	у0	z 0	
1	x1	x1	z1	
2	x2	y2	z2	
3	x3	уЗ	z3	
4	x4	y4	z 4	
5	x5	у5	z5	
6	хб	уб	z 6	
• • •	•••	• • •	• • •	



# Indexed Face Set

- Used in many formats
- Storage
  - Vertex: position
  - Face: vertex indices
- No *explicit* neighborhood info

Vertices				
v0	x0	уO	z 0	
v1	x1	x1	z1	
v2	x2	y2	z2	
v3	x3	yЗ	z3	
v4	x4	y4	z4	
v5	x5	у5	z5	
v6	хб	уб	z6	
•••	••	••	••	
	•	•	•	

Tria	Triangles				
t0	v0	v1	v2		
t1	v0	v1	v3		
t2	v2	v4	v3		
t3	v5	v2	v6		
•••	••	••	••		
	•	•	•		

**ETH** zürich



## Indexed Face Set: Problems

- Information about neighbors is not explicit
  - Finding neighbors vertices/edges/faces is O(#V)
  - Local mesh modifications cost O(V)



 Breadth-first search costs O(k\*#V) where k = # found vertices

## **Neighborhood Relations**

All possible neighborhood relationships:

1. Vertex	- Vertex	VV
2. Vertex	- Edge	VE
3. Vertex	- Face	VF
4. Edge	- Vertex	EV
5. Edge	- Edge	EE
6. Edge	- Face	EF
7. Face	- Vertex	FV
8. Face	- Edge	FE
9. Face	- Face	FF



We'd like O(1) time for queries and local updates of these relationships



- Introduce orientation into data structure
  - Oriented edges







- Introduce orientation into data structure
  - Oriented edges





- Introduce orientation into data structure
  - Oriented edges
- Vertex
  - Position
  - 1 outgoing halfedge index
- Halfedge
  - 1 origin vertex index
  - 1 incident face index
  - 3 next, prev, twin halfedge indices
- Face
  - 1 adjacent halfedge index
- Easy traversal, full connectivity





- One-ring traversal
  - Start at vertex





- One-ring traversal
  - Start at vertex
  - Outgoing halfedge




- One-ring traversal
  - Start at vertex
  - Outgoing halfedge
  - Twin halfedge





- One-ring traversal
  - Start at vertex
  - Outgoing halfedge
  - Twin halfedge
  - Next halfedge





- One-ring traversal
  - Start at vertex
  - Outgoing halfedge
  - Twin halfedge
  - Next halfedge
  - Twin ...





• Pros:

(assuming bounded vertex valence)

- O(1) time for neighborhood relationship queries
- O(1) time and space for local modifications (edge collapse, vertex insertion...)
- Cons:
  - Heavy requires storing and managing extra pointers
  - Not as trivial as Indexed Face Set for rendering with OpenGL / Vertex Buffer Objects



### Halfedge Libraries

#### CGAL

- www.cgal.org
- Computational geometry
- OpenMesh
  - www.openmesh.org
  - Mesh processing
- We will not implement a half-edge data structure in the class. Instead we will work with Indexed Face Set and augment it to have fast queries.

Roi Poranne



#### Next time...

# Surface reconstruction



Roi Poranne



## Thank you



February 21, 2018